Towards Managing Complex Data Sharing Policies with the Min Mask Sketch

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What are data sharing policies?

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- A sharing policy is a set of expressions that describe how, when, and what data can be accessed.
- Examples:
 - ACL's
 - IAM (Amazon Web Services)
 - Friend-based sharing
 - BitTorrent / Distributed data networks
 - Advertisements

What are simple data sharing policies?

A single expression describes how to share the data.





What are **complex** data sharing policies?

Multiple expressions describe how to share the data.

Sharing Policy ID(s)	Data
1	Record 1
3	Record 2
2	Record 3
1, 3	Record 4
1, 2, 3	Record 5

Example: Weather Company X



Example: Health Tracker Pro



Example Data Set

time	heart_rate	blood_sugar	body_temp
2016-02-20 04:05:06	71	95	98.6
2016-02-20 04:05:09	72	96	98.7
2016-02-20 04:05:09	72	94	98.7

•

2016-02-21 11:14:40	115	125	99.3
2016-02-21 11:14:43	115	124	99.5
2016-02-21 11:14:46	116	124	99.6

Example Data Set with Sharing Policies

time	heart_rate	blood_sugar	body_temp	high_hr	low_bs	high_bt
2016-02-20 04:05:06	71	95	98.6	0	1	0
2016-02-20 04:05:09	72	96	98.7	0	1	0
2016-02-20 04:05:09	72	94	98.7	0	1	0
		•				

2016-02-21 11:14:40	115	125	99.3	1	0	1
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How can we store this policy metadata more efficiently?

Probabilistic Data Structures

- Sacrifice a small amount of accuracy in exchange for space efficiency.
- Can answer queries about the data without needing to store the entire data set.
- Examples
 - Bloom Filter
 - Count Min Sketch



Bloom Filter

- Probabilistic data structure that is used to test whether an element is a member of a data set.
- Uses an array of bits and a collection of hash functions.
- Conceived by Burton Howard Bloom in 1970.

How Does it Work?

• Initialization:

Bloom Filter

How Does it Work?

- Initialization:
 - Set each bit in the array to 0.
 - Create k hash functions using technique from Kirsch et. al 2005

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- Insert an element, X.
- Let *k* = 3

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- Insert an element, X.
- Let *k* = 3
 - $\circ \quad h_1(X) = 7$

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- Insert an element, X.
- Let *k* = 3
 - \circ h₁(X) = 7
 - $\circ h_2(X) = 2$

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- Insert an element, X.
- Let *k* = 3
 - \circ h₁(X) = 7
 - $\circ h_2(X) = 2$
 - $\circ h_{3}(X) = 11$

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

- Insert an element, X.
- Let *k* = 3
 - \circ h₁(X) = 7
 - $\circ h_2(X) = 2$
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- Each hash value corresponds to an index in the array of bits.

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

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 - $\circ h_3(X) = 11$
- Each hash value corresponds to an index in the array of bits.
- For each index calculated above, set the associated bit to 1.

	Bloom Filter													
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

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	Bloom Filter													
0	0	0	0	0	0	0	1	0	0	0	0	0	0	0
	7													

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0	0	1	0	0	0	0	1	0	0	0	0	0	0	0
	1	1 2	I	1	1	1				1	1	1		

- Insert an element, X.
- Let *k* = 3
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• Query an element, W.

	Bloom Filter													
0	0	1	0	0	0	0	1	0	0	0	1	0	0	0

- Query an element, W.
- Hash W using all *k* hash functions.

	Bloom Filter													
0	0	1	0	0	0	0	1	0	0	0	1	0	0	0

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- Hash W using all *k* hash functions.
 - h₁(W) = 5
 - $h_{2}(W) = 2$
 - \circ h₃(W) = 1

	Bloom Filter													
0	0	1	0	0	0	0	1	0	0	0	1	0	0	0

- Query an element, W.
- Hash W using all *k* hash functions.
 - h₁(W) = 5
 - $h_{2}(W) = 2$
 - $\circ h_{3}(W) = 1$



- If all bits are 1, W is said to exist in the set.
- If all bits are **not** 1, W is said to not exist in the set.



Bloom Filter: False Positives

• Hash collisions can result in false positives.

	Bloom Filter													
0	0	1	0	0	0	0	1	0	0	0	1	0	0	0

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Bloom Filter: False Positives

- Hash collisions can result in false positives.
- $h_2(W)$ collided with $h_2(X)$
- If the result of all *k* hash functions collided with any other element, all the bits would be 1, even though W is not an element in the data set.



Bloom Filter: False Negatives are Not Possible

• If an element exists in the data set, the Bloom Filter query will always return true.

	Bloom Filter													
0	0	1	0	0	0	0	1	0	0	0	1	0	0	0

Count-min Sketch

- Like a Bloom Filter but uses an array of counters instead of an array of bits.
- Used to determine an element's frequency within a data set.
- Cormode et al. (2005)



Count-min Sketch: Inserting

- When inserting an element, the element's primary key is hashed using all *d* hash functions.
- The counter value at each index is then incremented.



Count-min Sketch: Querying

- When querying an element, the element's primary key is hashed using all *d* hash functions.
- The minimum counter value at each index is returned as the estimated frequency for the element.



Count-min Sketch: Frequency Estimates

- The frequency can be overestimated due to hash collisions.
- The frequency cannot be underestimated.


Count-min Sketch: Parameters

- Sketch is sized according to the desired quality.
- The frequency estimate is bounded by an additive factor of ε with probability c.
- **c** and **c** are chosen by the developer.



Min Mask Sketch

- Like a Count-min Sketch but uses an array of bit strings instead of an array of counters.
- Used to determine an element's sharing policy information within a data set.
- This paper.



What Does the Bit String Represent?

• Each position in the bit string represents a possible expression to evaluate in order to share or restrict data.

Expression 1	heart_rate > 114	
Expression 4	random() < 0.167	
Expression 8	LIMIT = 10	

00101001

What Does the Bit String Represent?

- Each position in the bit string represents a possible expression to evaluate in order to share or restrict data.
- If a bit at a particular position is set to 1, that expression is *active*

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00101001 **Expression 4** is active

What Does the Bit String Represent?

- Each position in the bit string represents a possible expression to evaluate in order to share or restrict data.
- If a bit at a particular position is set to 1, that expression is *active*.
- If a bit at a particular position is set to 0, that expression is *inactive*.



• The new element is hashed based on its primary key (x) using the *d* different hash functions.



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Min Mask Sketch: Querying

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- The bit string with the minimum number of 1's (active expressions) is returned as the estimated sharing policy bit string.



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- <u>https://github.com/oudalab/mms</u>



Workflow



Usage: Creating an Empty Min Mask Sketch

CREATE EXTENSION mms;

```
CREATE TABLE example (
    example_sketch mms
);
```

INSERT INTO example VALUES(mms());

Usage: Inserting an Element



Usage: Querying the Min Mask Sketch

SELECT mms_get_mask(example_sketch, "abc"::text)
 FROM example;

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- For **1** GB of data, The simple approach would require **187.5** MB.
- This results in the Min Mask Sketch providing a **187.49** MB reduction in storage cost for this example.

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- Cannot deactivate an expression (move from a 1 to a 0).
- When policies cluster together, the mms can become inefficient.

Future Directions

- Expanding the Min Mask Sketch to store types of metadata other than sharing policy information.
- Rigorous study of the performance characteristics of the Min Mask Sketch.
- Comparison with other solutions to handling sharing policies.

References

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Images Used

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Thank You!

Policy Log Approach

• What if the data sharing policies tend to cluster together?
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time	heart_rate	blood_sugar	body_temp	high_hr	low_bs	hide_bt
2016-02-20 04:05:06	71	95	98.6	0	1	0
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Policy Log Approach

- A log of the data sharing policies and when they change would be a better approach.
- This approach requires more space as a function of the policy changes.

key	high_hr	low_bs	high_bt
2016-02-20 04:05:06	0	1	0
2016-02-21 11:14:40	1	0	1

Min Mask Sketch vs. Policy Log

- In the context of the Health Tracker Pro example.
- Min Mask Sketch parameters:

• **c** = 99%

